

An Introduction to Functional Reactive Programming

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Reactive Programming

- **Reactive Program**: continually interacts with its environment in a **timely** manner.
- Examples: video games, mp3 players, robot controllers, aeroplane control systems ...
- Contrast with:
 - **Transformational Programs**, e.g. a compiler
 - **Interactive Programs**, e.g. accessing a database

What type of program?

Greeting

```
greeting = do putStrLn "What is your first name?"  
             n1 ← getLine  
             putStrLn "And what is your family name?"  
             n2 ← getLine  
             putStrLn ("Hello " ++ n1 ++ " " ++ n2)
```

Insertion Sort

```
isort :: Ord a ⇒ [a] → [a]  
isort []           = []  
isort (x : xs) = insert x (isort xs)  
  
insert :: Ord a ⇒ a → [a] → [a]  
insert x []           = [x]  
insert x (a : as) | x > a = a : insert x as  
                   | x ≤ a = x : a : as
```

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Greeting

Interactive

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              n1 ← getLine
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Insertion Sort

Transformational

```
isort :: Ord a => [a] -> [a]
isort []      = []
isort (x : xs) = insert x (isort xs)

insert :: Ord a => a -> [a] -> [a]
insert x []      = [x]
insert x (a : as) | x > a = a : insert x as
                  | x ≤ a = x : a : as
```

Functional Reactive Programming (FRP)

- FRP languages are domain-specific languages (the domain being reactive programming)
- Key characteristic: **inherent notion of time**
- Usually embedded in a host language (often Haskell)
- Also useful for modelling and simulation

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- The original idea of FRP was to provide a **continuous-time abstraction** to the FRP programmer...

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- The original idea of FRP was to provide a **continuous-time abstraction** to the FRP programmer...
- ...while automating the discretisation necessary for implementation.

Signals and Events

- FRP is based around **time-varying values** called signals (or behaviours):

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- There are also **instantaneous occurrences** called events.
- One way to represent events is as *Maybe* types within signals:

$$\text{Signal } (\text{Maybe } a)$$

Signal Functions

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- Some languages go further and **only** provide functions on signals as a first-class abstraction.

Signal Functions

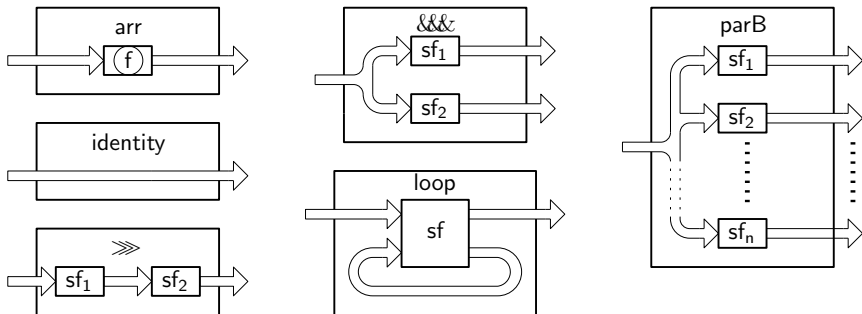
- FRP languages keep signals abstract, providing several signals, and functions on signals, as primitives.
- This has several advantages, e.g.
 - enforcing **causality**
 - optimisation opportunities
- Some languages go further and **only** provide functions on signals as a first-class abstraction.
- These are called **signal functions**:

$$SF\ a\ b \approx Signal\ a \rightarrow Signal\ b$$

Yampa: An FRP Language

- A DSL embedded in Haskell
- No signals, only signal functions
- Pretends to have continuous time
- Has been used for a variety of applications: video games, sound synthesis, robot simulators, GUIs, virtual reality, visual tracking, animal monitoring. . .

Yampa Routing Combinators



$arr \quad :: (a \rightarrow b) \rightarrow SF \ a \ b$

$identity \quad :: SF \ a \ a$

$(\gg\gg) \quad :: SF \ a \ b \rightarrow SF \ b \ c \rightarrow SF \ a \ c$

$(\&\&\&) \quad :: SF \ a \ b \rightarrow SF \ a \ c \rightarrow SF \ a \ (b, c)$

$parB \quad :: [SF \ a \ b] \rightarrow SF \ a \ [b]$

$loop \quad :: SF \ (a, c) \ (b, c) \rightarrow SF \ a \ b$

Some Yampa Primitives

Events

```
data Event a = NoEvent | Event a
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tag :: Event a → b → Event b
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Time-Dependent Primitives

```
integral :: Num a ⇒ SF a a
```

```
delay :: Time → a → SF a a
```

```
edge :: SF Bool (Event ())
```

```
switch :: SF a (b, Event e) → (e → SF a b) → SF a b
```

Examples

Example Yampa Code

```
localTime :: SF a Time
```

```
localTime = arr (const 1) >>> integral
```

```
after :: Time → SF a (Event ())
```

```
after t = localTime >>> arr (≥ t) >>> edge
```

```
integral :: Num x ⇒ x → SF x x
```

```
integral x = integral >>> arr (+x)
```

```
switchWhen :: SF a b → SF b (Event e) → (e → SF a b) → SF a b
```

```
switchWhen sf sfe = switch (sf >>> (identity &&& sfe))
```

Arrow Notation

- Yampa uses a special **do** notation (from the Arrow framework)

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Pure Code $(f :: a \rightarrow x)$

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 $\lambda (a, b) \rightarrow$   
  let  $x = f\ a$   
       $y = g\ (b, x)$   
  in  $h\ (x, y, b)$ 
```


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```

Monadic Code $(f :: a \rightarrow m\ x)$

```
 $\lambda (a, b) \rightarrow$  do  
   $x \leftarrow f\ a$   
   $y \leftarrow g\ (b, x)$   
   $h\ (x, y, b)$ 
```

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```

Arrow Code $(f :: SF\ a\ x)$

```
proc  $(a, b) \rightarrow$  do  
   $x \leftarrow f \multimap a$   
   $y \leftarrow g \multimap (b, x)$   
   $h \multimap (x, y, b)$ 
```

Bouncing Balls

See accompanying code. . .

Yampa Implementation

The SF data type (simplified)

```
data SF a b  $\approx$  SF (DTime  $\rightarrow$  a  $\rightarrow$  (SF a b, b))
```

(*DTime* is the amount of time passed since the previous sample.)

Summary

- FRP languages are domain-specific languages for reactive programming.
- Their key characteristic is an implicit notion of time.
- If you want to learn more about Yampa, I'd recommend Henrik Nilsson's recent mini-course:
<http://www.cs.nott.ac.uk/~nhn/ITU-FRP2010/>